

# Type Theory

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(they / them)

Strathclyde

Quantinuum

$$\sigma, \tau ::= \perp$$
$$| \sigma \rightarrow \tau$$
$$s, t ::= \lambda x. t$$
$$| \underline{e}$$
$$e, f ::= f s$$
$$| x$$
$$| t : \tau$$

$$\boxed{\tau \ni t}$$

$$\boxed{e \in \sigma}$$

$$\frac{x \in \sigma \vdash \tau \ni t}{\sigma \rightarrow \tau \ni \lambda x.t}$$

$$\sigma \rightarrow \tau \ni \lambda x.t$$

$$\frac{e \in \sigma \quad \sigma \equiv \tau}{\tau \ni e}$$

$$\tau \ni e$$

$$\frac{f \in \sigma \rightarrow \tau \quad \sigma \ni s}{f s \in \tau}$$

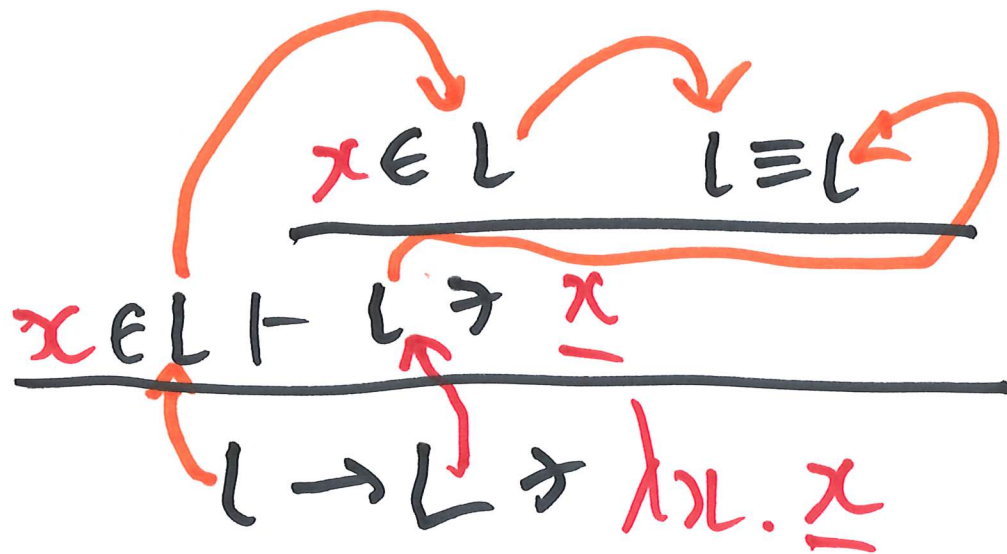
$$f s \in \tau$$

$$\frac{\tau \ni t}{t : \tau \in \tau}$$

$$t : \tau \in \tau$$

$$(\lambda x.t : \sigma \rightarrow \tau) s \rightsquigarrow t[s:\sigma] : \tau$$

$$\underline{t : \tau} \rightsquigarrow t$$



$(lx. ly. xyy) st$

$\rightsquigarrow$   
 $(ly. syy)t$

$\rightsquigarrow$   
 $stt$

$s, t, S, T ::= \star$   
|  $\Pi x:S. T$   
|  $\lambda x. t$   
|  $e$

$e, f ::= x$   
|  $f s$   
|  $t:T$

$$\{\star \Rightarrow T\} \boxed{T \Rightarrow \boxed{t} / \{\}} \{ \}$$

$$\{ \} \boxed{e \in S} / \{\star \Rightarrow S\} \{ \}$$

$$\overline{\star \Rightarrow \star}$$

$$\frac{f \in \Pi x:S. T \quad S \Rightarrow S}{f s \in T [s:S]}$$

$$\frac{\star \Rightarrow S \quad x \in S \vdash \star \Rightarrow T}{\star \Rightarrow \Pi x:S. T}$$

$$\frac{\star \Rightarrow T \quad T \Rightarrow t}{t : T \in T}$$

$$\star \Rightarrow \Pi x:S. T$$

$$x \in S \vdash T \Rightarrow t$$

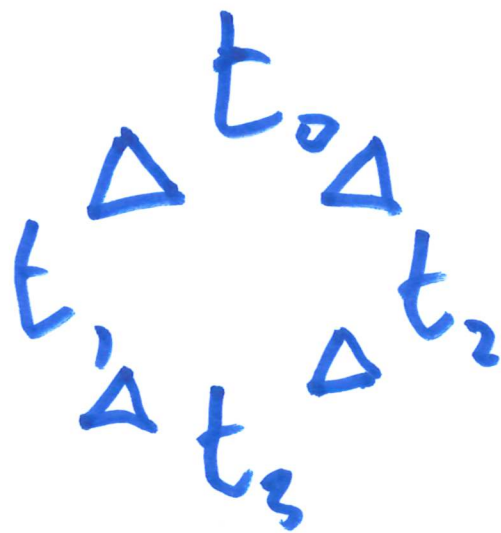
$$\Pi x:S. T \Rightarrow \lambda x. t$$

$$\frac{e \in S \quad S \equiv T}{T \Rightarrow e}$$

$$\frac{T \rightsquigarrow T' \quad T' \Rightarrow t}{T \Rightarrow t}$$

$$\frac{e \in S \quad S \rightsquigarrow S'}{e \in S'}$$

$$\begin{array}{c}
 \underline{t:T} \rightsquigarrow t \\
 \in \Pi x:S'.T'' \\
 \hline
 (\lambda x.t : \Pi x:S.T) s \rightsquigarrow (t:T)[s:S] \\
 \in T'[s:S]
 \end{array}$$





☆:☆:☆

$\Pi x: (\lambda x. \underline{x} : \Pi x: \star. \star) \star$

$(\lambda y. \underline{y} : \Pi y: \star. \star) \star$

$$\begin{array}{r}
 s \in S \quad t \in T \cancel{[s:S]} \\
 \hline
 s, t \in \Sigma_{x:S.T} \\
 \\
 S \ni s \quad \uparrow [s:S] \ni t \\
 \hline
 \Sigma_{x:S.T} \ni s, t
 \end{array}$$

SAT

EDF

t : T Δ T'

x Δ x t : T Δ t' : T'

$\frac{x \Delta t \Delta t'}{\lambda x. t \Delta \lambda x. t'}$

$\frac{f \Delta f' \quad s \Delta s'}{f s \Delta f' s'}$

~~$\frac{\lambda x. T \Delta T'}{t \Delta t'}$~~   
 $t : T \Delta t'$

$\frac{x t \Delta t' s \Delta s' \quad \lambda x. T \Delta T' s \Delta s'}{(\lambda x. t : T \Delta t') s \Delta (t : T') [s' : s']}$

$$\Gamma \vdash T \Rightarrow t$$

$$\begin{matrix} \Delta^* & \Delta^* & \Delta \\ \Gamma' \vdash & T' \Rightarrow & t' \end{matrix}$$

$x \in S \Rightarrow x \in S'$

$\Gamma \vdash e \in S$	$S$
$\Gamma' \vdash e' \in S'$	$S'$

$$\forall \Gamma, T, t. \Gamma \vdash T \Rightarrow t \rightarrow \forall$$

$$\forall \Gamma, \Gamma', T, T', t, t'. \Gamma \vdash T \Rightarrow t \wedge \Gamma \Delta^* \Gamma' \wedge T \Delta^* T'$$

$$\wedge t \Delta t'$$

$$\rightarrow \Gamma' \vdash T' \Rightarrow t'$$

$$\forall \Gamma, \Gamma', e, e', S. \Gamma \vdash e \in S \wedge \Gamma \Delta^* \Gamma' \wedge e \Delta e'$$

$$\rightarrow \exists S'. S \Delta^* S' \wedge \Gamma' \vdash e' \in S'$$

need

$$T_2 \rightsquigarrow^* T_3$$

$$T_1 \triangleright^* T_3$$

$$\frac{T_0 \rightsquigarrow T_1 \quad T_1 \ni t}{T_0 \ni t}$$

$$T_0 \ni t$$

$$\begin{array}{l} \triangleright^* \\ T_2 \end{array} \quad \begin{array}{l} \triangleright \\ t' \end{array}$$

$$\begin{array}{l}
 e' \in S_2 \rightsquigarrow^* S_3 \\
 e \in S_0 \Delta^* \quad S_0 \rightsquigarrow S_1 \\
 \hline
 \end{array}$$

$$\begin{array}{l}
 e \in S_1 \\
 \nabla \Delta^* \\
 e' \in S_3
 \end{array}$$